
Formalization and Analysis of the Separation Minima for Aircraft in NAT

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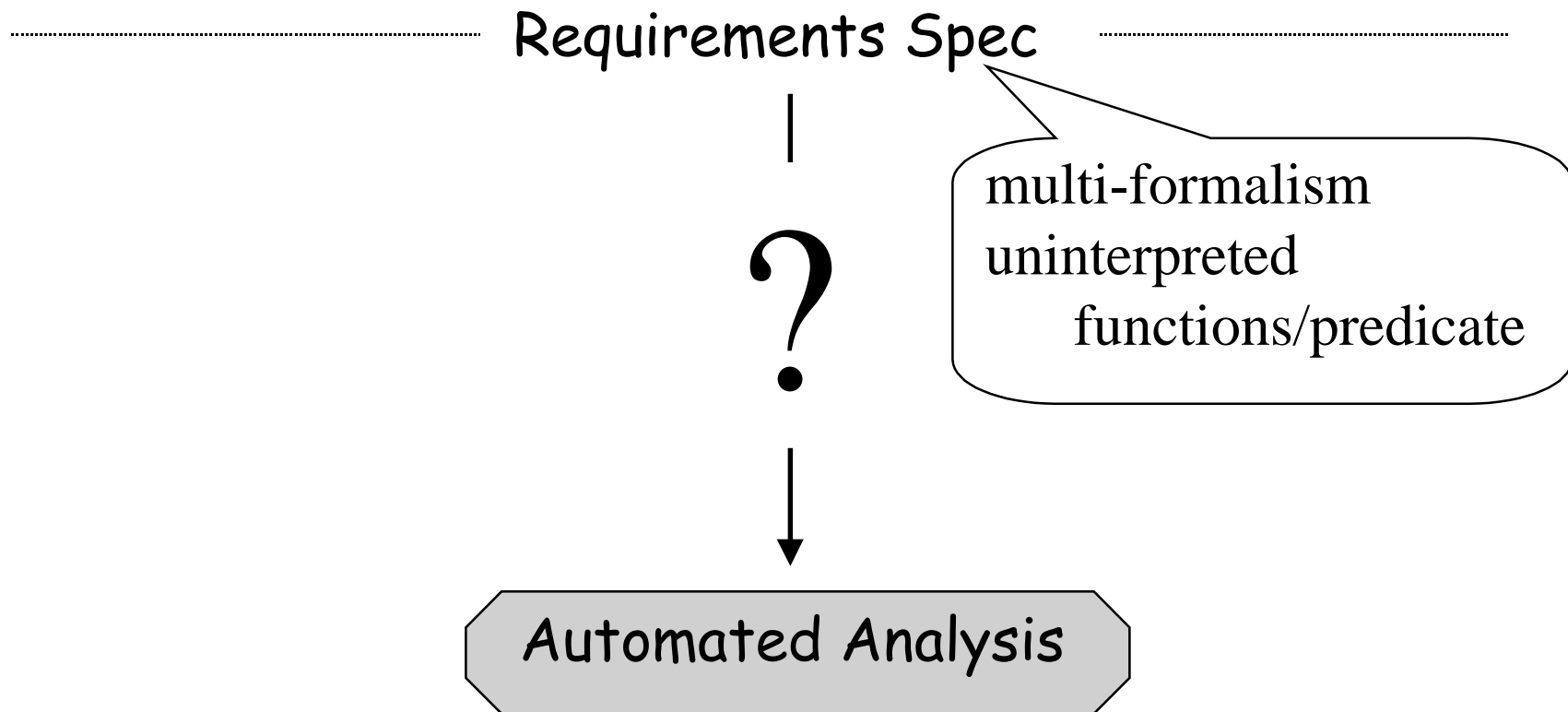
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Thesis: Formal Validation of System Specifications

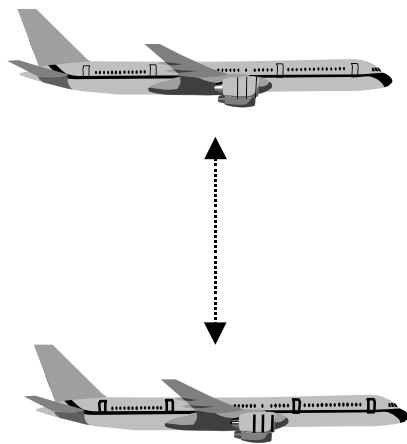


Outline

- system
 - formal description technique
 - analysis techniques
 - analysis results
 - future work
 - summary
-

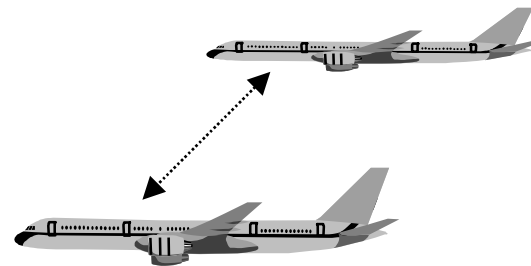
Separation Minima for NAT

rules for air traffic controllers



Vertical
Separation

Lateral Separation



Longitudinal Separation



Existing Specification

- complex decision logic
- stateless

Form of the specification:

- informal natural language (64 pgs)
 - thoroughly reviewed and in use
 - pseudocode interpretation (9 pgs)
-

Analysis

■ completeness

- are all combinations of inputs for two flights covered ?
- what are the cases covered by the default values ?

■ consistency

- is it ever possible that the specification has multiple possible outcomes for the same inputs ?
-

Goals

- both specification and analysis results had to be readable and reviewable by the domain expert
 - formal specification:
 - in a suitable notation
 - don't add details
 - analysis - return results:
 - in the terms given in the specification
-

Formal Specification

- combination of:
 - ASCII based predicate logic
 - » defined types, functions and predicates types
 - » uninterpreted types, functions and predicates
 - function and predicate tables
 - » slight variation of AND/OR tables
 - 15 tables; 16 defns; 47 uninterpreted functions and predicates
 - 18 pg document; formal spec = 300 lines
-

Example

:flight;

FlightLevel : flight -> num;

IsLevel: flight -> bool;

:typeOfAircraft := Turbojet | Supersonic | Other;

TypeOfAircraft : flight -> typeOfAircraft;

← OR →

↑ AND ↓				Default
	FlightLevel(A)	_ < 280	_ > 450	
	TypeOfAircraft(B)	_ = Turbojet	_ = Supersonic	
	IsLevel(A)	_ = T	.	
	InCruiseClimb(A)	.	_ = F	
	Separation (A,B)	1000	4000	2000

structure captures related elements in a row

Definitions of Completeness and Consistency of a Table

- completeness:
 - default cases
 - consistency:
 - no two columns with differing result values for the function overlap
 - with respect to possible values for row entries
-

Outline

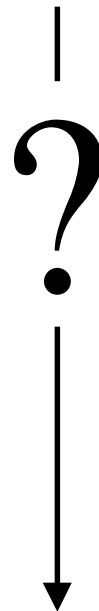
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Analysis :Context

Specifier

..... Requirements Spec

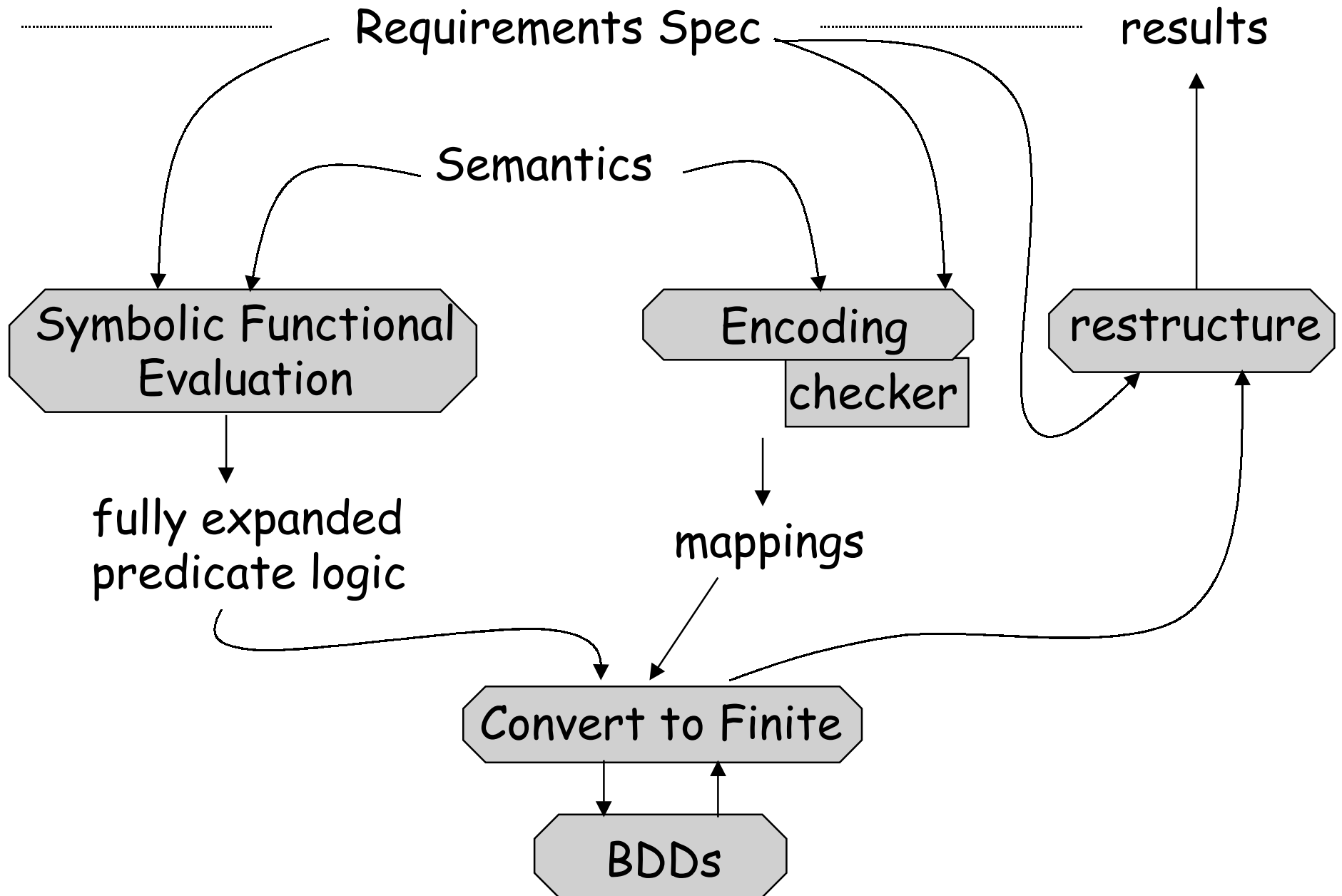
Heimdahl and Leveson
Heitmeyer, et al.
Hoover and Chen



simplifications
to make finite

Automated State Space Exploration Analysis
(Finite State Machines; BDDs)

Framework



Key features of this approach

- direct use of explicit defn of semantics
 - integrate notations
 - symbolic functional evaluation
 - encodings based on structure
 - reverse mappings and structuring of results
-

Analysis Results

- completeness analysis found:
 - missing assumptions "everyone knew about" (domain knowledge)
 - incorrect partitions
 - consistency analysis found:
 - three places where the requirements are inconsistent
 - symmetry analysis found:
 - assumptions about the uninterpreted terms
-

Environmental Assumptions

- predicate logic constraints

forall A:flight. NOT (IsLevel(A) AND InCruiseClimb(A));

- analysis:

- applied to existing items of the correct type
- existentially quantified out of result

- limited output

- allowed symmetry checking to be more successful

Future Work

- extending framework to do model checking of specifications with state, such as statecharts
 - aeronautical telecommunications network (ATN)
-

Summary

■ analysis technique:

- handles uninterpreted functions/predicates
- definition of the semantics used in analysis
 - » integrate multiple notations
- takes advantage of structure found in tables
- user can add environmental assumptions
- return results in terms and structure of specification



extend what could be done fully automatically

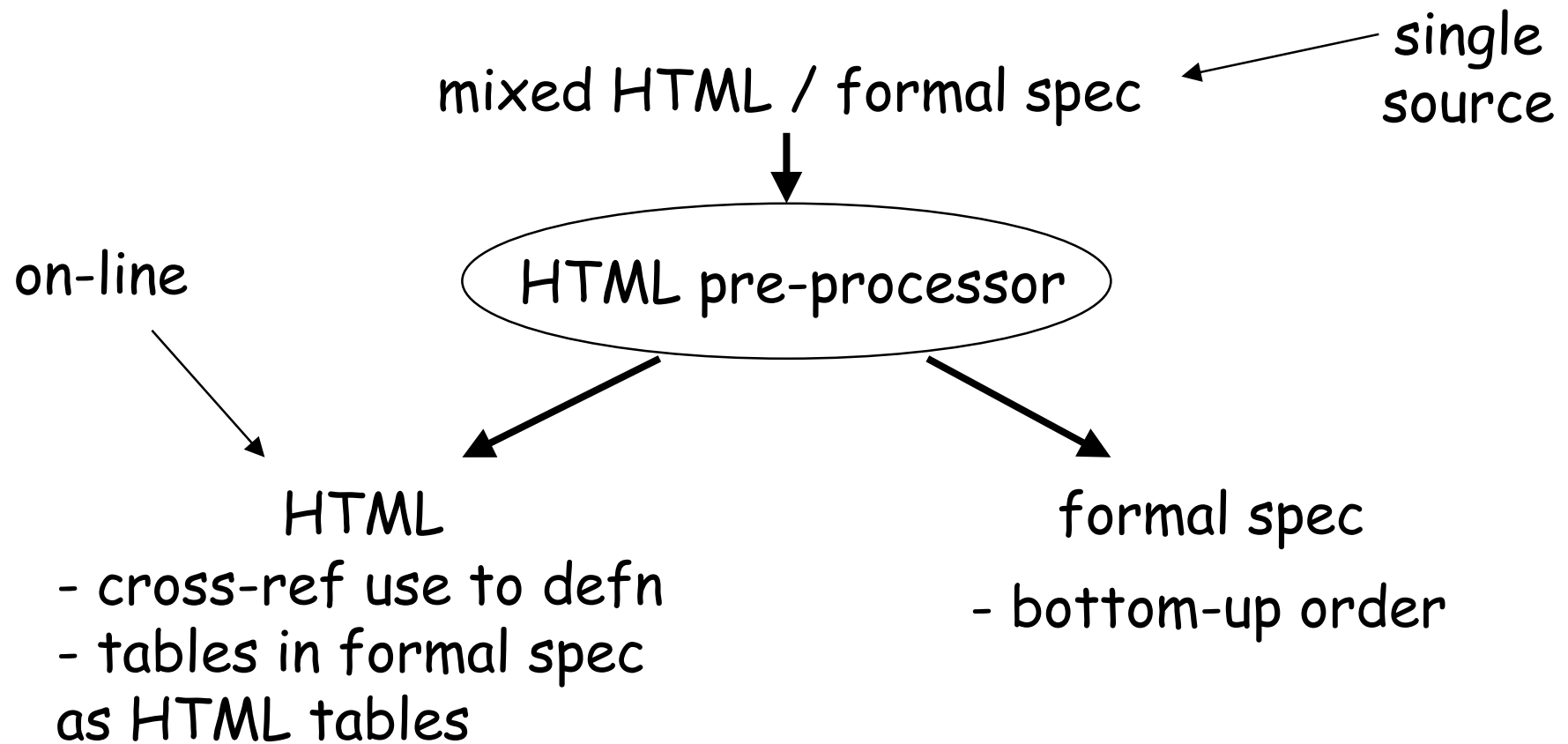
Thesis: Validation of System Specifications

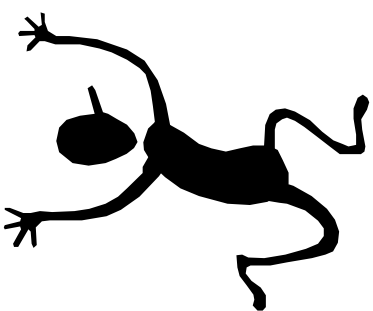
- multi-formalism specifications
- uninterpreted functions/predicates
- exploit structure in specifications
- let user control the analysis through their knowledge of the domain



extend range of what can be automatically analyzed

Presentation of the Specification





That's all

Example of Analysis Results

Advantages / Contributions

- use the explicit defn of semantics directly in analysis; also simulation, prototyping; analysis of semantics
 - general framework for:
 - multiple notations; multiple analysis techniques
 - non-formal methods person; formal methods expert
 - return results at correct level of abstraction
 - exploit inherent abstractions
-

Related Work

- builds on existing work by
 - Heimdahl and Leveson
 - Heitmeyer,
 - especially for definitions of compl/cons
 - PVS
 - similar approach but
 - » didn't have to add a construct to the language
 - » compl/cons not required by defn of semantics
 - » enumerating resultant cases
 - » took advantage of environment
-

Embedding

- had to get into a common form for analysis
 - S: ASCII notation based on higher order logic; open tool support
 - predicate logic parts were given directly in S
 - tables were embedded in S using a textual representation
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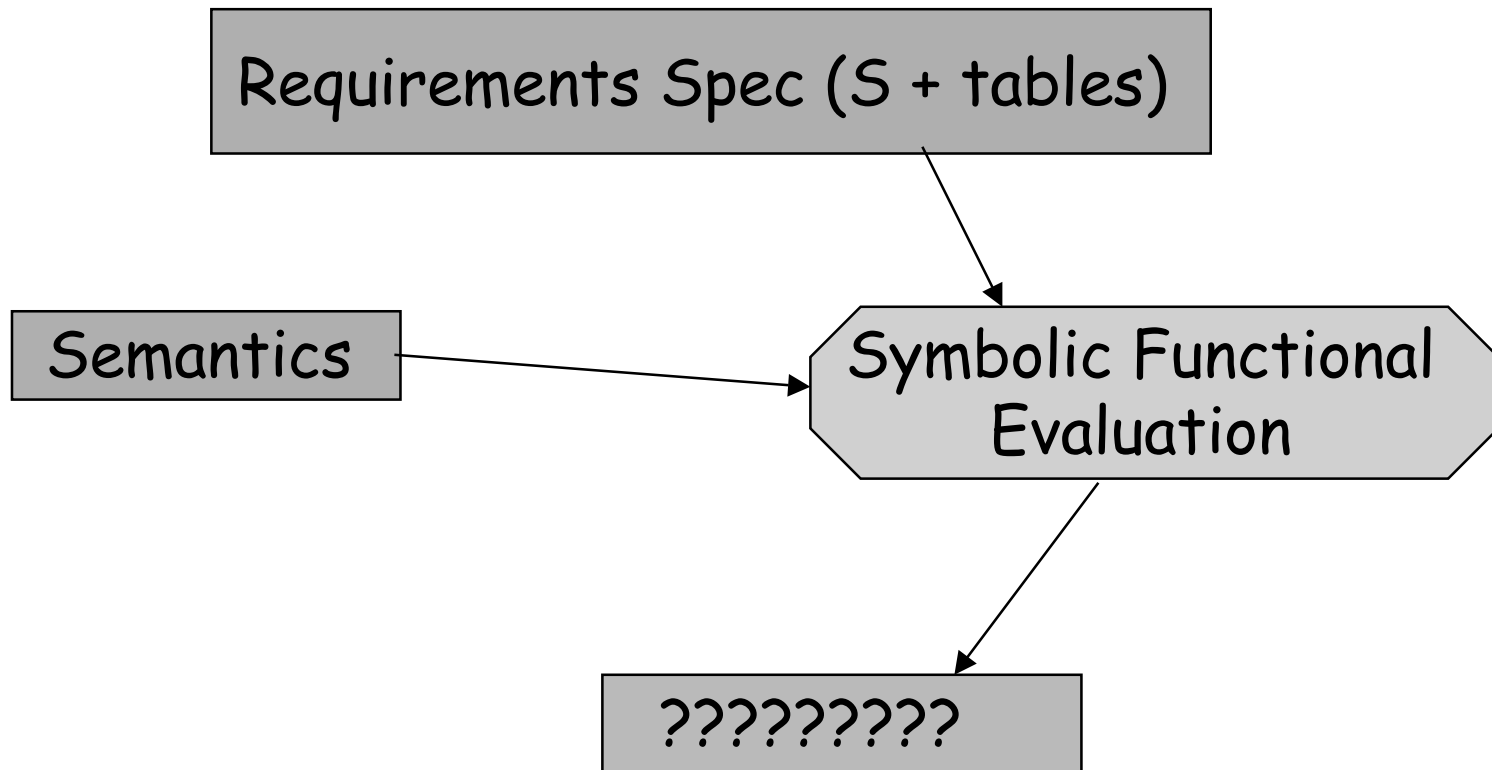
Semantics of the tables

- also written in S (given in Appendix)
 - keyword "Table" or "Predicate Table" is the semantic function
 - "executable" in that they can reduce the table into something in terms of the entries in the rows and Boolean connectives
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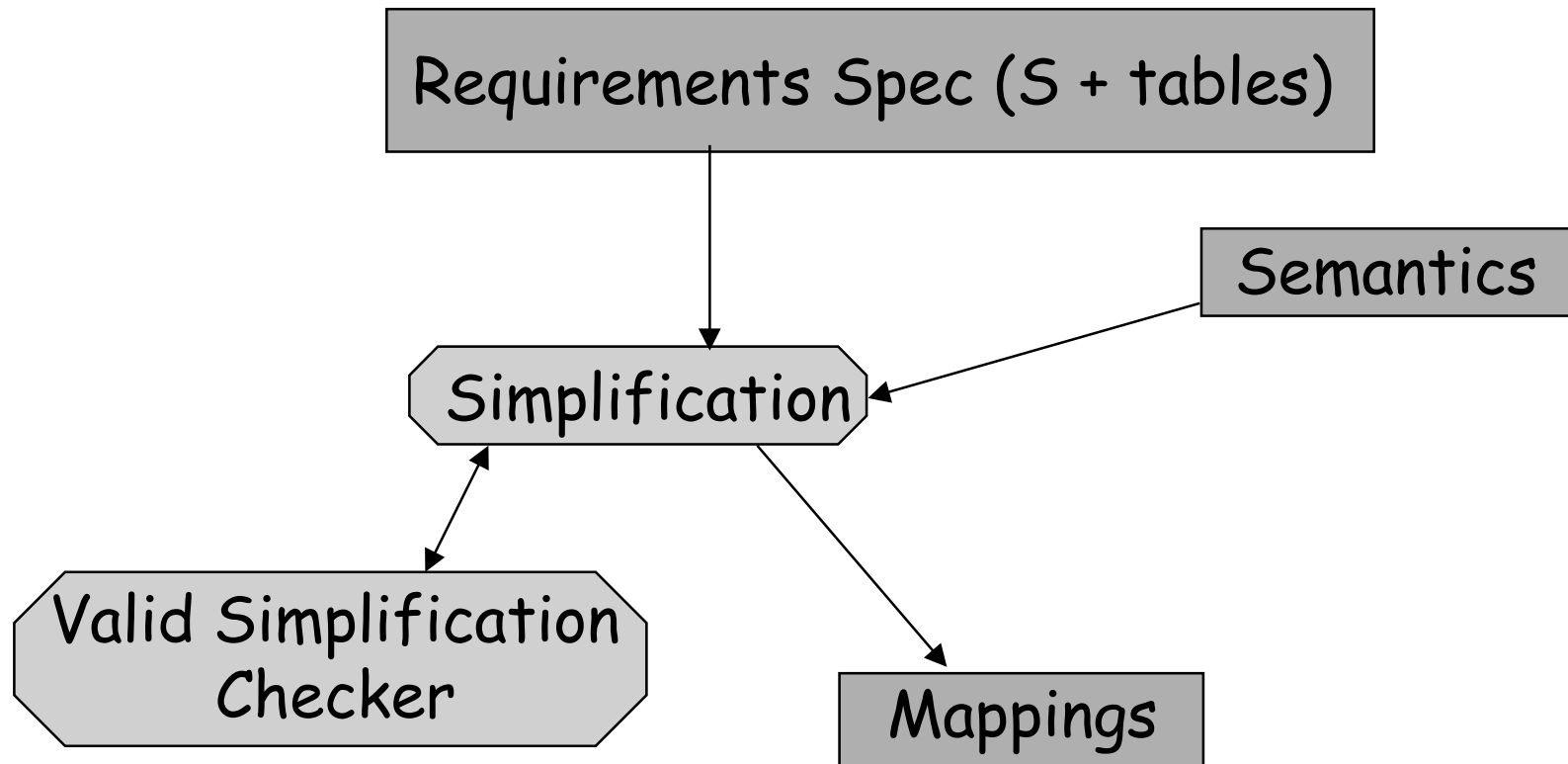
Thesis Statement

Having an explicit machine-readable operational semantics for a notation within a common framework provides a systematic way to exploit inherent abstractions to carry out state-space exploration analysis.

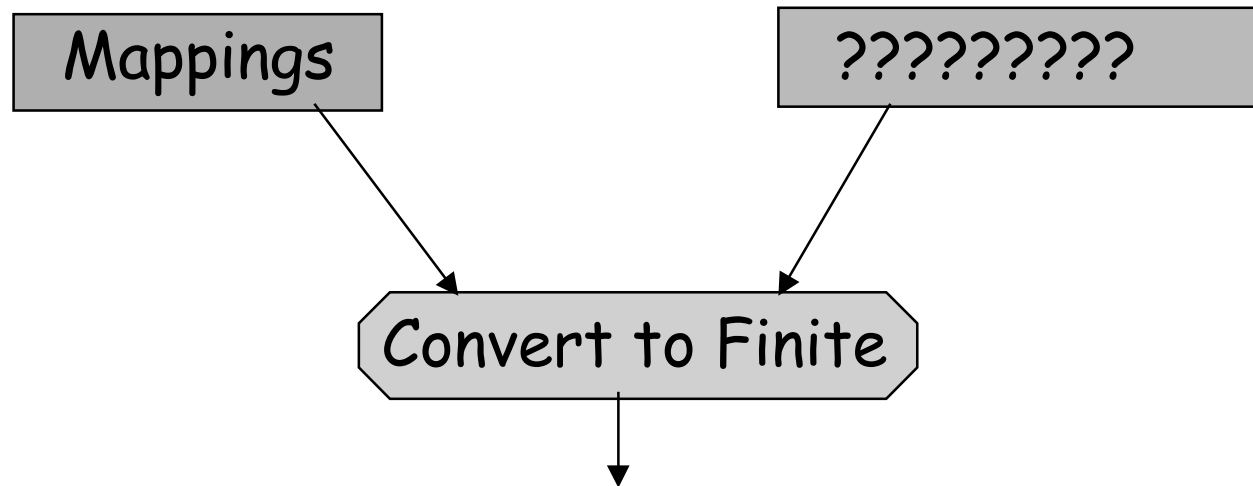
First Step: Dealing with Multiple Notations



Second Step: Determining Valid Abstractions

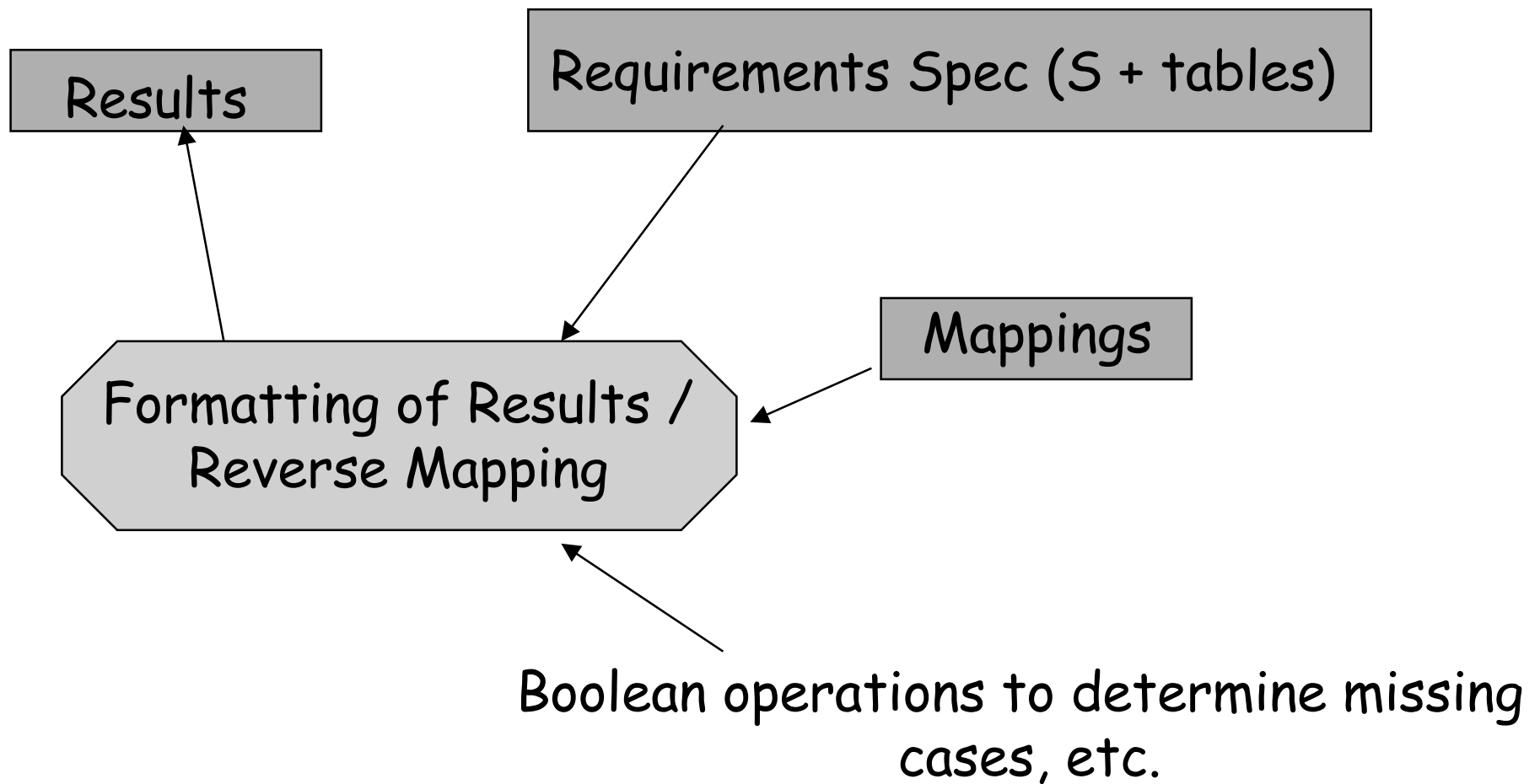


Third Step: Mapping to a Finite Domain and ???



Boolean operations to determine missing cases, etc.

Fourth Step: Returning Results



fix this picture

