

CPSC 421/501

Oct 11, 2024

- Characterization of
regular languages over $\{a\}$

- Myhill-Nerode point of
view of regular languages
over $\{a\}$ and beyond

↳ See handout:

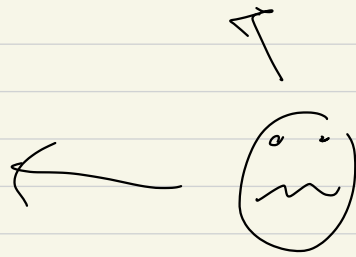
"Non-regular languages and
the Myhill-Nerode theorem"

Myhill-Nerode theorem

classifies all regular

languages, at least in

principle ---



$$L = \{ a^n \mid n \geq 2, n \in \mathbb{N},$$

and n and $n+2$
are prime numbers } }

Is L regular?

Easy:

L is regular iff there are

finitely many n s.t.

$n, n+2$ are prime.

iff the "twin prime

conjecture" is false

3, 5 twin primes

5, 7 " "

11, 13 " "

17, 19 " "

29, 31 " "

Solve this \Rightarrow good position, good salary

Languages over $\{a\} = \Sigma$

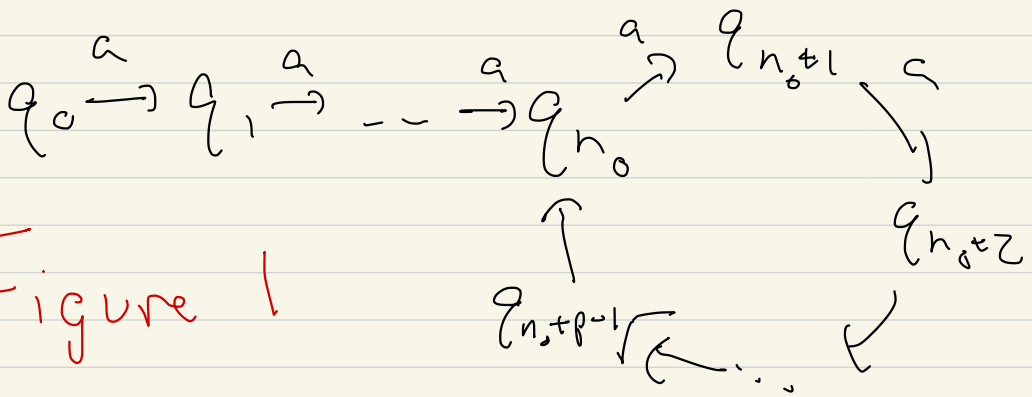
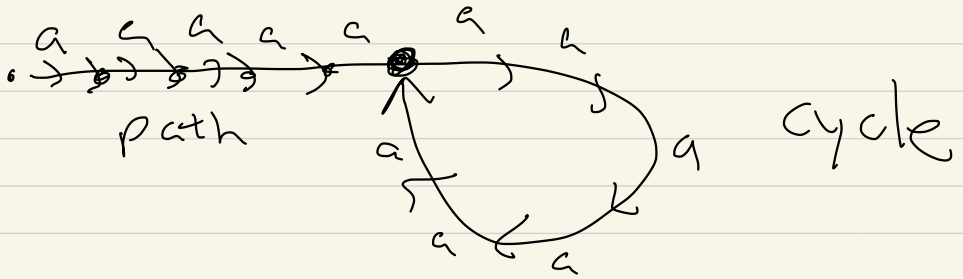


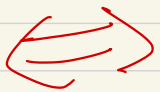
Figure 1

path length n_0
(n_0 directed edges)

cycle of
 p edges

$n_0 + p$ edges \Leftrightarrow $n_0 + p$ vertices
 \Leftrightarrow $n_0 + p$ states

If L is regular, then there
is a DFA with $n_0 + p$ states
that looks like Figure 1



Which strings reach

q_{n_0} : a^{n_0} , a^{n_0+p} , a^{n_0+2p} , ...

if q_{n_0} is accepting:

a^{n_0} , a^{n_0+p} , a^{n_0+2p} , ... $\in L$

and if q_{n_0} is rejecting, not
accepting, then a^{n_0} , a^{n_0+p} , ... $\notin L$

Similarly, for $i = 0, 1, \dots, p-1$

q_{n_0+i} get there via

$$a^{n_0+i}, a^{n_0+i+p}, a^{n_0+i+2p}, \dots$$

either all these $\in L$

OR " " $\notin L$

Upshot: If L is regular,
then for some n_0 and p

$$n_0 \geq 0, p \geq 1$$

$$n_0 \in \mathbb{Z}_{\geq 0}, p \in \mathbb{N} = \{1, 2, \dots\}$$

for all $i = 0, \dots, p-1,$

$a^{n+1}, a^{n+2}, a^{n+3}, \dots$

either all these $\in L$

OR " " $\notin L$

(*)

i.e.

$a^n \in L \Leftrightarrow a^{n+p} \in L$

for all $n \geq n_0$

(**)

Result: smallest p , then smallest $n_0 \rightarrow$

Define: If $L \subseteq \{a\}^*$,

then we say L is

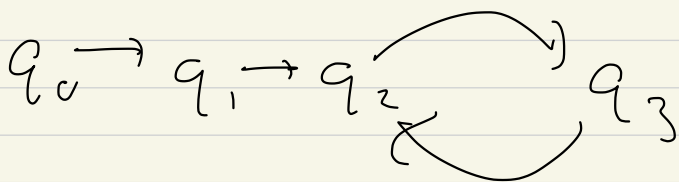
eventually p -periodic if

(*) or (**) hold

for some p .

If so, the smallest value
of p for which L is
eventually is called the
eventual period of L .

e.g. if



$$L = \left(\begin{array}{c} a^0 \\ \text{or} \\ \text{not} \end{array} \right), \left(\begin{array}{c} a^1 \\ \text{or} \\ \text{not} \end{array} \right), \left(\begin{array}{c} a^2, a^4, a^6, \dots \\ \text{or} \\ \text{not} \end{array} \right),$$

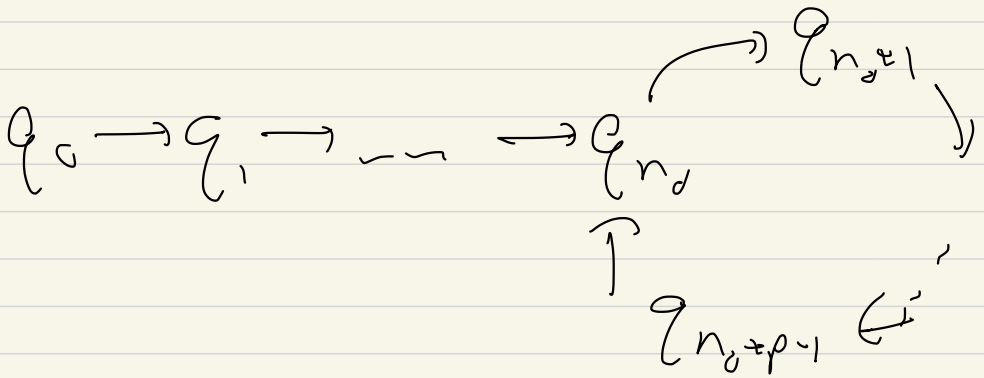
$$\text{and } \left(\begin{array}{c} a^3, a^5, a^7, \dots \\ \text{or} \\ \text{not} \end{array} \right)$$

Thm! If L is eventually periodic with eventual period p , and n_0 is the smallest non-neg integer

such that

$$\forall n \geq n_0, a^n \in L \Leftrightarrow a^{n+p} \in L$$

The DFA recognizing L with the fewest states has $n_0 + p$ states and look like

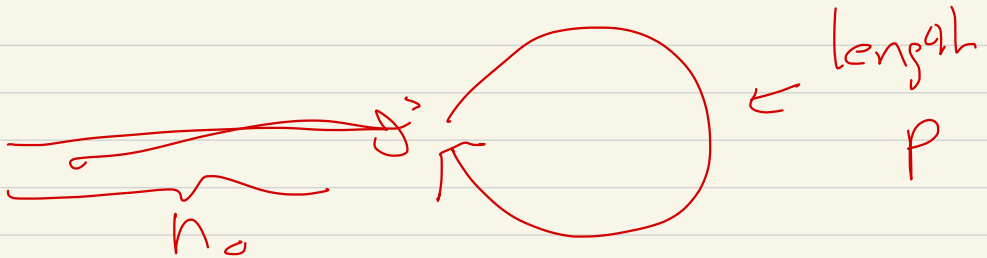


Example :

$$\{ a, a^4, a^9, a^{16}, a^{25}, \dots \}$$

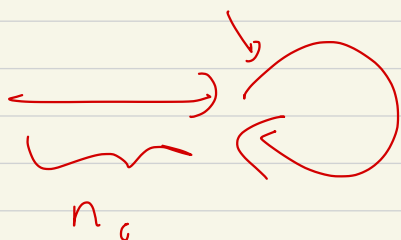
$$= \{ a^{n^2} \mid n \in \mathbb{N} \}$$

is not regular, since



$$(q+1)^2 - q^2 = 2q+1 \geq q+1$$

$a^{q^2+1}, a^{q^2+2}, \dots, a^{q^2+q}$
 not in L



p states on the cycle

$$\text{if } |a^{q^2+1}| \geq |a^{n_0}|$$

$$q^2+1 \geq n_0$$

$$q \geq \sqrt{n_0}$$

so all states in cycle reject

Similarly $n_1 < n_2 < n_3 < \dots$

sr,

$$L = \{ a^{n_1}, a^{n_2}, a^{n_3}, \dots \}$$

$$\text{and } (n_{i+1} - n_i) \xrightarrow{i \rightarrow \infty} \infty$$

then L is not regular

Exercise (Kevin Liu)

$$\text{say that } \left(n_1 = 3, n_2 = 4, 5 \mid n_3 = 6, n_4 = 7, 8 \right)$$

Could be a nice homework
or midterm problem - -

are the countably many
or uncountably "

What if $n_{i+1} = n_i + 1$

for i large

$L = \{ a^2, a^5, a^9, a^{16}, \dots \}$

eventual period = 1

$n \geq 9$

$a^n \in L \Leftrightarrow a^{n+1} \in L$

What about if L is
finite,

i.e. $\{a\}^* \not\subseteq L$ then

this includes $a^n, a^{n+1}, a^{n+2}, \dots$

some n large

period = 1

$\{ a^n \mid n, n+2 \text{ are primes} \}$