

CPSC 421/501

- Multi-tape \leq One-tape

- Multi-tape:

Function vs Decision

- Addition / ADDITION
- Multiplication / MULTIPLICATION

- Decide

A_{DFA} = Acceptance for DFA's

Recognize: A_{TM} = " " " TM

- Take questions \approx 10 min

- " " " on Wednesday

\approx 20 min

- Midterm: Friday, Room 6

- N.B.: Supplemental Practice
for the 2024 midterm now
has 7th question on countability.

- All HW good sample midterm
problems...

You can have 1 2-sided

Letter (8.5" x 11") sheet of

notes, handwritten

Computer

reduced

Level of detail

{ See solution to HW
Go to office hours

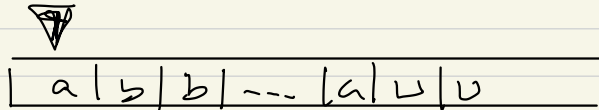
As we delve into multi-tape
Turing machines, and more
sophisticated algorithms ...

• formal description (i.e. $Q, \Sigma, \Gamma, \delta$)
to
"higher-level" ...

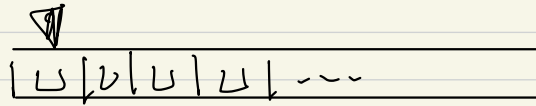
First abstraction: multi-tape
Turing machine



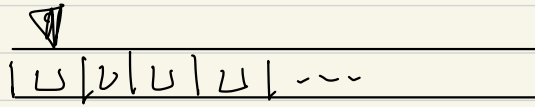
Q



tape 1



(input appears here)



tape 2

← tape 3

(some write #
of tapes, k,

$$\delta: Q \times \Gamma^k \rightarrow Q \times \Gamma^k \times \{L, R, S\}^k$$

$$\Gamma = \Sigma \cup \{U\} \cup \left\{ \begin{array}{l} \text{any others} \\ \text{you wish} \\ \text{to add} \end{array} \right\}$$

stay

finitely many

Why --

1	← carry	← line
138		← one line
+	238	← another line
376		← another line

four lines \leadsto 4-tape machine

ADDITION = $\left\{ \begin{array}{l} x \# y \# z \\ x, y, z \in \{0, 1, \dots, 9\}^* \\ \text{s.t. } x + y = z \end{array} \right\}$

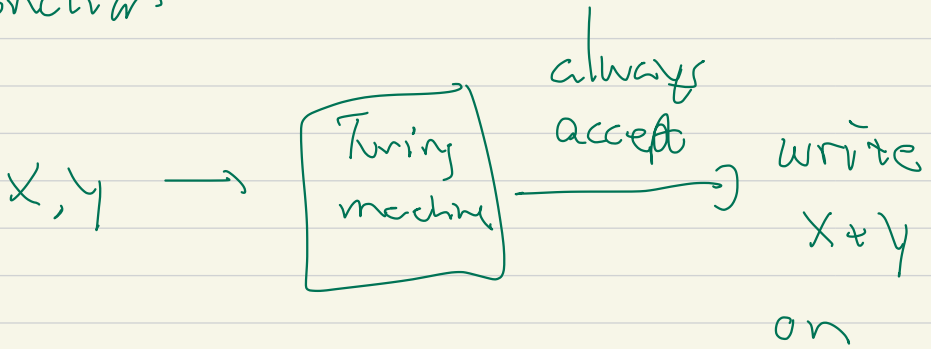
Given $x, y, z \Rightarrow$ yes $x + y = z$
 no $x + y \neq z$

addition! takes

12517 # 3916 uuu

returns their sum

function



Turing machines
that compute
functions

some
tape

ADDITION $\subset \Sigma^*$

$$\Sigma = \{c, 1, \dots, 9\} \cup \{\#\}$$

MULTIPLICATION / multiplication:

big {

$$\begin{array}{r} 3265 \\ \times 121 \\ \hline \end{array}$$

many lines {

$$\begin{array}{r} 3265 \\ + 65300 \\ \hline 326500 \end{array}$$

} finite #

OK

$$\begin{array}{r} 3265 \\ \times 2 \\ \hline 6530 \end{array}$$

End of New Content,

Brief Q's about practice

midterm. -

{ all maps $\mathbb{N} \rightarrow \{1, 2\}$ }

\updownarrow 1-1 bijection
to

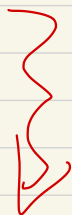
Power(\mathbb{N})

$f: \mathbb{N} \rightarrow \{1, 2\} \rightarrow \{ \text{yes, no} \}$

$1 \mapsto \text{yes}$
 $2 \mapsto \text{no}$

So given

$$f: \mathbb{N} \rightarrow \{\text{yes}, \text{no}\}$$



$$\{ n \in \mathbb{N} \mid f(n) = \text{yes} \}$$

$$\{ \cancel{1}, 2, 3, \cancel{4}, 5, \dots \}$$

↓ ↓ ↓ ↓ ...
no yes yes no

\tilde{f} :

$1 \mapsto \text{no}$

$2 \mapsto \text{yes}$

$3 \mapsto \text{no}$

$4 \mapsto \text{yes}$

$5 \mapsto \text{yes}$

Subset of \mathbb{N}
corresp to \tilde{f}

2

4

5

⋮
⋮
⋮

$\text{Power}(\mathbb{N}), \text{Power}(\Sigma^*)$

are uncountable

$S \xrightarrow{\text{injection}} T, |T| \geq |S|$

$A \xrightarrow{\text{surjection}} B, |A| \geq |B|$

So

$\mathbb{C} \xrightarrow{\text{surjection}} \text{Power}(\mathbb{N})$
(uncountable)
 \mathbb{C} must be uncountable

